Survivable Network Design for Group Connectivity in Bounded-Treewidth Graphs

Daniel Vaz

Parinya Chalermsook, Syamantak Das, Guy Even, Bundit Laekhanukit, MPI Informatics, Germany Aalto Univ., Finland IIIT Delhi, India Tel-Aviv Univ., Israel MPI Informatics, Germany

Group SNDP



Group SNDP

- Input:
 - Graph G, edge/vertex costs
 - Groups S_i with connectivity demands k_i
 - Root r
- Goal: Find a min-cost graph that
 - Contains k_i disjoint paths from r to $v_i \in S_i$

Known Results

- $k_i = 1$: Group Steiner Tree
 - $O(\log^3 n)$ -approx*, $O(\log^2 n)$ -hardness [GKR'98], [HK'03]
 - Bounded Treewidth: $O(\log^2 n)$ -approx [CDLV17]
- $k_i \leq 2$ • $ilde{O}(\log^4 n)$ -approx*

[GKR'10]

- * Cannot be improved beyond $O(\log n)$
 - Both use tree embedding, which has distortion $\Omega{\log n}$

Known Results

[KKN'12]

- Groups of size 1: SNDP
 - 2-approx [Jain'00]
- In general: Label-Cover hard
 - $2^{\log^{1-\epsilon} n}$ hardness
 - $\Omega(n^{\delta})$ hardness under Sliding Scale Conjecture

Our Results

- Group SNDP:
 - $O(\log^2 n)$ -approx in $n^{f(k,tw(G))}$
- SNDP:
 - Exact algorithm in time $n2^{f(k,tw(G))}$
- Both results for vertex weights / connectivity

Main Idea

- Develop a DP without group constraints
- Turn the DP into an LP and add the group constraints
- New LP represents a variant of Group Steiner Tree

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Fine-Grained Complexity and Algorithms



Ramamohan Paturi

UC San Diego Foundations of Finegrained Complexity





Amir Abboud

IBM Almaden

Hardness in P





Danupon Nanongkai

KTH

Dynamic graphs: algorithms, conditional lower bounds, and complexity classes